There are three basic constructs in set theory:

| Cartesian product | $S \times T$ |
| :--- | :--- |
| Power set | $\mathbb{P}(S)$ |
|  |  |
| Comprehension 2 | $\{x \mid P\}$ |

where $S$ and $T$ are sets, $x$ is a variable and $P$ is a predicate.

Jean-Reymond Abrial (ETH-zurich)

|  | ${ }_{\text {Baxic Constructs }}^{\substack{\text { Extenions }}}$ |
| :---: | :---: |
| Set Comprehension |  |



These axioms are defined by equivalences.

| Left Part | Right Part |
| :---: | :---: |
| $E \mapsto F \in S \times T$ | $E \in S \wedge F \in T$ |
| $S \in \mathbb{P}(T)$ | $\forall x \cdot(x \in S \Rightarrow x \in T)$ <br> $(\mathrm{x}$ is not free in S and T$)$ |
|  |  |
| $E \in\{x \mid P\}$ | $[x:=E] P$ <br> $(\mathrm{x}$ is not free in E$)$ |


| Left Part | Right Part |
| :---: | :---: |
| $S \subseteq T$ | $S \in \mathbb{P}(T)$ |
| $S=T$ | $S \subseteq T \wedge T \subseteq S$ |

The first rule is just a syntactic extension
The second rule is the Extensionality Axiom


Elementary Set Operator Memberships

| $E \in S \cup T$ | $E \in S \vee E \in T$ |
| :--- | :--- | :--- |
| $E \in S \cap T$ | $E \in S \wedge E \in T$ |
| $E \in S \backslash T$ | $E \in S \wedge E \notin T$ |
| $E \in\{a, \ldots, b\}$ | $E=a \vee \ldots \vee E=b$ |
| $E \in \varnothing$ | $\perp$ |
| Seander |  |



| Generalized Union | union $(S)$ |
| :--- | :--- |
| Union Quantifier | $\cup x \cdot(P \mid T)$ |
| Generalized Intersection | $\operatorname{inter}(S)$ |
| Intersection Quantifier | $\cap x \cdot(P \mid T)$ |



Basic Const
Extensions
Generalized Intersection


| $E \in \operatorname{union}(S)$ | $\exists s \cdot s \in S \wedge E \in s$ <br> $(\mathrm{~s}$ is not free in S and E$)$ |
| :--- | :--- |
| $E \in(\bigcup x \cdot P \mid T)$ | $\exists x \cdot P \wedge E \in T$ <br> $(\mathrm{x}$ is not free in E$)$ |
| $E \in \operatorname{inter}(S)$ | $\forall s \cdot s \in S \Rightarrow E \in s$ <br> $(\mathrm{~s}$ is not free in S and E$)$ |
| $E \in(\bigcap x \cdot P \mid T)$ | $\forall x \cdot P \Rightarrow E \in T$ <br> $(\mathrm{x}$ is not free in E$)$ |



| union $(S)$ |
| :--- |
| $\bigcup x \cdot P \mid T$ |
| $\operatorname{inter}(S)$ |
| $\cap x \cdot P \mid T$ |



| Binary relations | $S \leftrightarrow T$ |
| :--- | :--- |
| Domain | $\operatorname{dom}(r)$ |
| Range | $\operatorname{ran}(r)$ |
| Converse | $r^{-1}$ |


$r \in A \leftrightarrow B$

$\operatorname{ran}(r)=\{b 1, b 2, b 4, b 6\}$



| Domain restriction | $S \triangleleft r$ |
| :--- | :--- |
| Range restriction | $r \triangleright T$ |
| Domain subtraction | $S \notin r$ |
| Range subtraction | $r \triangleright T$ |

Summary of the Mathematical Notation


The Range Restriction Operator

$F \triangleright\{b 2, b 4\}$

$\{a 3, a 7\} \notin F$

$F \triangleright\{b 2, b 4\}$


| Image | $r[w]$ |
| :--- | :--- |
| Composition | $p ; q$ |
| Overriding | $p \not q q$ |
| Identity | $i d(S)$ |





$$
\begin{aligned}
& r^{-1-1}=r \\
& \operatorname{dom}\left(r^{-1}\right)=\operatorname{ran}(r) \\
& (S \triangleleft r)^{-1}=r^{-1} \triangleright S \\
& (p ; q)^{-1}=q^{-1} ; p^{-1} \\
& (p ; q) ; r=q ;(p ; r) \\
& (p ; q)[w]=q[p[w]] \\
& p ;(q \cup r)=(p ; q) \cup(p ; r) \\
& r[a \cup b]=r[a] \cup r[b]
\end{aligned}
$$

$$
\begin{array}{ll}
\text { Given a relation } r \text { such that } r \in S \leftrightarrow S \\
& \\
r=r^{-1} & \forall x, y \cdot x \in S \wedge y \in S \Rightarrow(x \mapsto y \in r \Leftrightarrow y \mapsto x \in r) \\
r \cap r^{-1}=\varnothing & \forall x, y \cdot x \mapsto y \in r \Rightarrow y \mapsto x \notin r \\
r \cap r^{-1} \subseteq \operatorname{id}(S) & \forall x, y \cdot x \mapsto y \in r \wedge y \mapsto x \in r \Rightarrow x=y \\
\operatorname{id}(S) \subseteq r & \forall x \cdot x \in S \Rightarrow x \mapsto x \in r \\
r \cap \operatorname{id}(S)=\varnothing & \forall x, y \cdot x \mapsto y \in r \Rightarrow x \neq y \\
r ; r \subseteq r & \forall x, y, z \cdot x \mapsto y \in r \wedge y \mapsto z \in r \Rightarrow x \mapsto z \in r
\end{array}
$$

Set-theoretic statements are far more readable than predicate calculus statements

More classical Results

Given a relation $r$ such that $r \in S \leftrightarrow S$

| $r=r^{-1}$ | $r$ is symmetric |
| :--- | :--- |
| $r \cap r^{-1}=\varnothing$ | $r$ is asymmetric |
| $r \cap r^{-1} \subseteq \operatorname{id}(S)$ | $r$ is antisymmetric |
| $\operatorname{id}(S) \subseteq r$ | $r$ is reflexive |
| $r \cap \operatorname{id}(S)=\varnothing$ | $r$ is irreflexive |
| $r ; r \subseteq r$ | $r$ is transitive |

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|  | (extensions $\begin{gathered}\text { Basic Constructs } \\ \text { Exter }\end{gathered}$ |
| :---: | :---: |
| Function Operators (1) |  |


| Partial functions | $S \rightarrow T$ |
| :--- | :--- |
| Total functions | $S \rightarrow T$ |
| Partial injections | $S \leftrightarrow T$ |
| Total injections | $S \mapsto T$ |

A Partial Function F from a Set A to a Set B

$F \in A \rightarrow B$

$F \in A \rightarrow B$

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$F \in A \longmapsto B$


$F \in A \Longrightarrow B$

| Partial surjections | $S \rightarrow T$ |
| :--- | :--- |
| Total surjections | $S \rightarrow T$ |
| Bijections | $S \mapsto T$ |



$F \in A \rightarrow B$

| Left Part | Right Part |
| :---: | :---: |
| $f \in S \rightarrow T$ | $f \in S \rightarrow T \wedge T=\operatorname{ran}(f)$ |
| $f \in S \rightarrow T$ | $f \in S \rightarrow T \wedge T=\operatorname{ran}(f)$ |
| $f \in S \rightarrow T$ | $f \in S \rightarrow T \wedge f \in S \rightarrow T$ |



| $S \rightarrow T$ | $S \rightarrow T$ |
| :--- | :--- |
| $S \rightarrow T$ | $S \rightarrow T$ |
| $S \multimap T$ | $S \leftrightarrow T$ |
| $S \multimap T$ |  |



